# Geometric Data Analysis 

## Clustering

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Fall 2020

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(1) Clustering

- Vector spaces
- Improved k-means
- Arbitrary (non-vector) metric spaces
- Improvements and Evaluation


## Outline

## (1) Clustering

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## Clustering

## Definition (k clusters)

Given $n$ objects, and $k>1$, partition the objects into $k$ subsets (clusters) so as to optimize some objective function.

- Objects in the same cluster are more "similar" (or closer) to each other than to those in other clusters.
- Possible criteria: minimizing the total distance among all cluster points, minimizing the distance of cluster points to some center, etc.
- Variations: $k$ is unknown and computed, e.g., by the Silhouette method. Capacitated/balanced: $k$ given, clusters of equal cardinality.


## Good Clustering, with centers



## Approaches

- hierarchical (agglomerative): each point initializes a cluster, merge closest pair (define distance of clusters) until stopping criterion, e.g., predetermined number of clusters, or clusters with points too far apart (Gromos).
- point-assignment: given some initial clusters, assign points to "best" cluster; cluster represented by "centroid" (which may not belong to input). Example: k-means.
(Ullman et al.:Mining Massive datasets)


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## k-means: Objective function

Typical ambient space is $\mathbb{R}^{d}$ but can generalize to metric space $\mathcal{Z}$.

## Minimization function

In any metric space over points/vectors $\mathcal{Z}$ with distance/metric function d , let the dataset be $X=\left\{x_{1}, \ldots, x_{n}\right\} \subseteq \mathcal{X} \subseteq \mathcal{Z}, k>1$. Given centroids $C \subset \mathcal{Z}$, let

$$
\mathrm{d}\left(x_{i}, C\right)=\min _{c \in C} \mathrm{~d}\left(x_{i}, c\right) .
$$

Consider vector $v(C)=\left(d\left(x_{1}, C\right), \ldots, d\left(x_{n}, C\right)\right)$. The $k$-means objective is:

$$
\min _{C \subseteq \mathcal{Z},|C|=k}\|v(C)\|_{2}^{2}=\sum_{i=1}^{n} \mathrm{~d}\left(x_{i}, C\right)^{2}
$$

The $k$-means objective is NP-hard, but for the $\ell_{2}$ metric, Lloyd's algorithm converges quickly to a local minimum.

## Variations

## Various minimizations

Recall $X=\left\{x_{i}\right\}, v(C)=\left(\mathrm{d}\left(x_{1}, C\right), \ldots, \mathrm{d}\left(x_{n}, C\right)\right), C \subset \mathcal{Z}$ are centroids. For $\mathrm{d}(\cdot)$ denoting $\ell_{2}$ distance, the $k$-means objective is:

$$
\min _{C \subseteq \mathcal{Z},|\mathcal{C}|=k}\|v(C)\|_{2}^{2}=\sum_{i=1}^{n} \mathrm{~d}\left(x_{i}, C\right)^{2}
$$

Other standard objectives:
-- $k$-median: $\min _{C \subseteq \mathcal{Z},|C|=k}\|v(C)\|_{1}$,
-- $k$-medoid: $\min _{C \subseteq x,|C|=k}\|v(C)\|_{1}$.
-- $k$-center: $\min _{C \subseteq x,|C|=k}\|v(C)\|_{\infty}$.

## Lloyd's

## Algorithm

Initialize $k$ centers randomly (or using some strategy).
(1) Assignment: Assign each object to its nearest center.
(2) Update: Calculate mean $\frac{1}{T} \sum_{i=1}^{T} \overrightarrow{v_{i}}$ of each cluster, make it new center.

Repeat the two steps until there is no change in the assignments.

## Properties

- Each distance calculation $=O(d)$ because vectors in $\mathbb{R}^{d}$.
- Assignment $=O(n k d)$, Update $=O(n d)$,
- \#iterations unknown, in practice $\ll n$.
- Converges to local minimum in Euclidean space (depends on initialization)


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## Inverted Quantized k-means (IQ-means)

## Reverse assignment

- Fixed Data-structure for points (Dataset in memory)
- Centroids are queries; range search of increasing radius
- Resolve overlapping balls; consider "left-overs".

(Avrithis-Anagnostopoulos-Kalantidis-E,ICCV' 15)


## Assignment by Range search

## Reverse approach (ANN)

- Index $n$ points, but only once for entire algorithm.
- At each iteration, for each centroid c, range/ball queries centered at c.
- Mark assigned points: move at end of bucket, or flag them.
- Increase radii by $\times 2$, start with min(dist between centers)/2, until all points assigned, or most ranges/balls do not assign a new point.
- For a given radius, if a point lies in $\geq 2$ balls, compare its distances to the respective centroids, assign to closest centroid.
- At end: for every unassigned point, compare its distances to all centroids
- Standard method: $n$ ANN queries, each $k^{\rho}$, hence $O\left(n k^{\rho}\right)$.
- Inverse: $k$ queries, each $n^{\rho}+$ OutputSize $=O(n)$; stores entire datase $\dagger$probabilistic analysis


## Inverted Quantized k-means (IQ-means)

## The algorithm

- inverse assignment (above): faster than update!
- quantization on dynamic 2d-grid (via learning) (Avrithis:ICCV'13)
- dynamic estimation of overlap hence of $k$ (Avrithis-Kalantidis,ECCV' 12)



## Dynamic IQ-means (k=9)



## Dynamic IQ-means (k=7)



## Dynamic IQ-means (k=7)



## Experiments

- (Avrithis,Kalantidis,Anagnostopoulos,E:ICCV' 15) http://github.com/iavr/iqm
- Comparison against
-- AKM: Approximate $k$-means (Philbin et al. CVPR’07)
-- RR: Ranked Retrieval (Broder et al. Web Search \& Data Mining' 14)
-- standard $k$-means
- Speed: IQ-means fastest
- Accuracy: IQ-M on par with dedicated methods, worse than (approx) $k$-means
- clustering of 100 M images, on a single machine, in $<1$ hour. Best method for a couple of years.


## Performance

Distortion vs total time for 20 iterations on $10^{6}$ images (SIFT1M):



Left: varying number of clusters $k$. Right: increasing number of points $n$.

## Mining example

500K Paris + 100Mil YahooFlicker images


Accurate cluster despite large dataset: Paris ground truth depicted in red outline, the rest are images closest to the red ones.

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## k-medoids

Goal: Handle any distance metric; $k$-means only if consistent with mean.
k-medoids (PAM is simplest algorithm) use centroids that belong to the dataset:

## Definition (Medoid)

The medoid of a set is the object of the set that minimizes total dissimilarity (distance) to all other objects in the set.

Objective function (cf. above): Minimize sum of distances to point's centroid.

## vs k-means

-- k-means tends to select convex spherical clusters; $k$-medoids less so.
-- $k$-means is more sensitive to noisy data and outliers.
-- k-means is faster and easier to implement.
(Kaufman-Rousseeuw ${ }^{\text {'87 }}$

## Partitioning Around Medoids (PAM)

Initialize $k$ centroids randomly.
(1) Assignment: Assign each object to nearest centroid; compute objective
(2) Update:
for each centroid $m$ do for each non-centroid $t$ do

Swap $m$ and $t$, compute new objective function value.
end for

## end for

Keep configuration (centroids) with min objective value.
Repeat steps 1 and 2 until there is no change of configuration (centroids).

Let distance calculation $=O\left(d^{\prime}\right)$. Update of Objective $=O\left((n-k) d^{\prime}\right)$, if $2 n d$ best centroid known. Hence, update $=O\left((n-k)^{2} k d^{\prime}\right) \sim n^{2}$.

## Accelerating updates

Two faster updates, which may however lose accuracy compared to PAM. Recall that after every swap we compute $J$ in $O\left((n-k) d^{\prime}\right)$.

## 1. Improved Update

Instead of swapping centroid $m$ with every point $t$, swap $m$ only with every non-centroid in same cluster as $m$.

Complexity: $n-k$ iterations instead of $k(n-k)$, hence update $=O\left((n-k)^{2} d^{\prime}\right)$

## 2. Update à la Lloyd's

For every cluster: (i) compute its medoid $t$, (ii) swap its current centroid $m$ with $t$.
The medoid $t$ minimizes $\sum_{i \in C} d(i, t)$ over all possible objects $t$ in cluster $C$. Computed in $O\left(a^{2} d^{\prime}\right)$, assuming clusters have $a \simeq n / k$ items.
Total Complexity $=O\left(\left(k a^{2}+k(n-k)\right) d^{\prime}\right)=O\left(\left(n^{2} / k+n k\right) d^{\prime}\right)=O\left(n^{2} d^{\prime}\right)$ (Park-Jun’09).

## Initial






## PAM 1



## PAM 2





## PAM 3





## PAM 4



## PAM 5



## PAM 6




## Final





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## Clustering Large Applications (CLARA)

General Idea: run entire algorithm with sample of size $n^{\prime} \ll n$. Use $s$ samples drawn independently, return best clustering.

Overall algorithm:

$$
\text { for } i=1, \ldots, s \text { do }
$$

apply PAM on a random (uniform) sample of size $n^{\prime}$ assign $n$ points to $k$ computed centroids calculate the total cost of the partitioning
end for
return best partitioning

Experimental results recommend: $s=5, n^{\prime}=40+2 k$.

## CLARA based on RANdomised Search (CLARANS)

- Update: swap m's with $t$ 's, for some randomly selected $(m, t)$ 's only.
- Picking random $Q \subset\{1, \ldots, k\} \times\{1, \ldots, n-k\}$, $s$ times.

Select $k$ centroids by some initialization method.
for $i=1, \ldots, s$ do
Cluster $n-k$ points to $k$ centroids by some assignment method. Randomly select set $Q$ of $|Q|$ pairs $(m, t),|Q|<k(n-k)$.
for $(m, t) \in Q$ do
Swap $m$ with $t$; compute new objective value.
end for
Keep centroids with minimum objective value over $|Q|$ choices.
end for
Output centroids yielding minimum objective value over $s$ candidates.
Experiments recommend: $s=2,|Q|=\max \{0.12 \cdot k(n-k), 250\}$.
(Ng-Han:IEEE Tran.Know.Data Eng’02, Theodoridis et al.:Patt.Recogn.,ch.14)

## Improve Initialisation 1: Spread-out

## initialization++ : k-means++ / k-medoids++:

(1) Choose a centroid uniformly at random; $t \leftarrow 1$.
(2) $\forall$ non-centroid point $i=1, \ldots, n-t$, let $D(i) \leftarrow$ min distance to some centroid, among $\dagger$ chosen centroids.
(3) Choose new centroid: $r$ chosen with probability proportional to $D(r)^{2}$ :

$$
\operatorname{prob}[\text { choose } r]=D(r)^{2} / \sum_{i=1}^{n-t} D(i)^{2}
$$

Let $t \leftarrow t+1$.
(4) Go to (2) until $t=k=$ given \#centroids.

Expected approximation ratio $=O(\log k)\left(\right.$ Arthur-Vassilvitskii:SODA $\left.{ }^{\prime} 07\right)$ Similar algo for 2-approx of $k$-center (NP-hard prob)

## Improve Initialisation 2: Concentrate

Select centroids close to dataset's center of mass (and to each other) as follows.
(1) Calculate symmetric $n \times n$ distance matrix of all objects, i.e. all distances $d_{i j}$ from every object $i=1, \ldots, n$ to every other object $j=1, \ldots, n, i \neq j$.
(2) For object $i$ compute

$$
v_{i}=\sum_{j=1}^{n} \frac{d_{i j}}{\sum_{t=1}^{n} d_{j t}}, \quad i=1, \ldots, n .
$$

(3) Return the $k$ objects with $k$ smallest $v_{i}$ values.

Algorithm proposed in (Park-Jun’09).

## Evaluation: Silhouette

-- For $1 \leq i \leq n, a(i)=$ average distance of $i$ to other objects in same cluster.
-- Let $b(i)=$ average distance of $i$ to objects in next best (neighbor) cluster, i.e. cluster of $2 n d$ closest centroid.

## Silhouette of Object $i$

$$
s(i)=\frac{b(i)-a(i)}{\max \{a(i), b(i)\}}=\left\{\begin{array}{ll}
1-a(i) / b(i), & \text { if } a(i)<b(i) \\
0, & \text { if } a(i)=b(i) \\
b(i) / a(i)-1, & \text { if } a(i)>b(i)
\end{array}\right\} \in[-1,1]
$$

## Interpret silhouette

$s(i) \rightarrow 1$ : i seems correctly assigned to its cluster;
$s(i) \simeq 0$ : borderline assignment (but not worth to change);
$s(i) \rightarrow-1$ : $i$ would be better if assigned to next best cluster.

## Silhouette: Cluster and clustering

## Specific clusters

-- Evaluate a cluster: Compute average $s(i)$ over all $i$ in some cluster.
-- If $k$ is too large or too small, some clusters shall display much smaller silhouettes than the rest.
-- Silhouette plots are used to improve $k$ : try different $k$ 's and see if clusters have roughly equal silhouettes.

## Overall Clustering

Overall Silhouette coefficient = average $s(i)$ over $i=1, \ldots, n$.
High if well clustered, low may indicate bad $k$ (or existance of outlier points).

