

## One-dimensional $l_1$ cost function minimization

**Proposition:** Let  $x_1, x_2, \dots, x_n \in R$ , with  $x_1 < x_2 < \dots < x_n$ . Consider the quantity

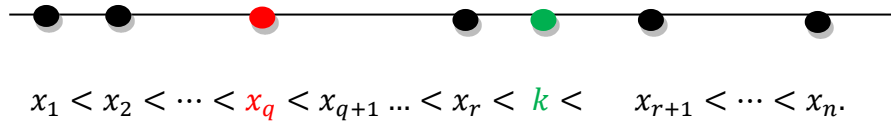
$$A = \sum_{i=1}^n |x_i - \mu|.$$

Then,  $A$  is minimized when  $\mu$  is chosen as the median of  $x_1, x_2, \dots, x_n$ , i.e.

$$\mu = \text{med}(x_1, x_2, \dots, x_n).$$

**Proof:**

(a) Let  $n$  be **odd** and  $x_q \equiv \mu = \text{med}(x_1, x_2, \dots, x_n)$ . Consider a number  $k > x_q$  so that  $x_r \leq k < x_{r+1}$ , with  $r > q$  (see the following figure).



Let  $k = x_q + \lambda$ , ( $\lambda > 0$ ).

Also, let

$$A_1 = \sum_{i=1}^n |x_i - x_q|,$$

and

$$A_2 = \sum_{i=1}^n |x_i - k|.$$

We consider the following three cases

(i)  $x_i \leq x_q$ : In this case it is  $|x_i - x_q| = x_q - x_i$  and  $|x_i - k| = k - x_i = x_q - x_i + \lambda$

(ii)  $x_q < x_i \leq x_r (\leq k)$ : In this case it is  $|x_i - x_q| = x_i - x_q$  and  $|x_i - k| = k - x_i = x_q - x_i + \lambda$

(iii)  $x_r (\leq k) < x_i$ : In this case it is  $|x_i - x_q| = x_i - x_q$  and  $|x_i - k| = x_i - k = x_i - x_q - \lambda$

It is

$$\begin{aligned} A_1 &= \sum_{i=1}^q |x_i - x_q| + \sum_{i=q+1}^r |x_i - x_q| + \sum_{i=r+1}^n |x_i - x_q| \\ &= \sum_{i=1}^q (x_q - x_i) + \sum_{i=q+1}^r (x_i - x_q) + \sum_{i=r+1}^n (x_i - x_q) \end{aligned}$$

and

$$\begin{aligned}
A_2 &= \sum_{i=1}^q |x_i - k| + \sum_{i=q+1}^r |x_i - k| + \sum_{i=r+1}^n |x_i - k| \\
&= \sum_{i=1}^q (x_q - x_i + \lambda) + \sum_{i=q+1}^r (x_q - x_i + \lambda) + \sum_{i=r+1}^n (x_i - x_q - \lambda) \\
&= \sum_{i=1}^q (x_q - x_i) + \sum_{i=1}^q \lambda + \sum_{i=q+1}^r (x_q - x_i) + \sum_{i=q+1}^r \lambda + \sum_{i=r+1}^n (x_i - x_q) \\
&\quad - \sum_{i=r+1}^n \lambda
\end{aligned}$$

Then, it is

$$\begin{aligned}
\Lambda &= A_2 - A_1 = \sum_{i=1}^q \lambda + \sum_{i=q+1}^r \lambda - \sum_{i=r+1}^n \lambda + \sum_{i=q+1}^r (x_q - x_i) - \sum_{i=q+1}^r (x_i - x_q) \\
&= \sum_{i=1}^r \lambda - \sum_{i=r+1}^n \lambda + 2 \sum_{i=q+1}^r (x_q - x_i) = (2r - n)\lambda - 2 \sum_{i=q+1}^r (x_i - x_q)
\end{aligned}$$

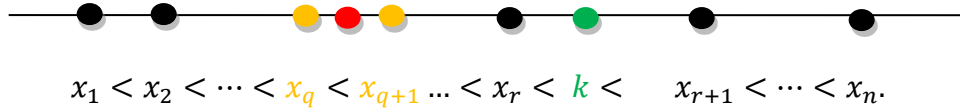
For  $i = q + 1, \dots, r$ , it is  $x_i - x_q \leq x_r - x_q \leq k - x_q = \lambda$ .

Therefore (taking also into account that  $q = \frac{n+1}{2}$ )

$$\Lambda = A_2 - A_1 \geq (2r - n)\lambda - 2 \sum_{i=q+1}^r \lambda = (2r - n)\lambda - 2(r - q)\lambda = (2q - n)\lambda = \left(2 \frac{n+1}{2} - n\right)\lambda = \lambda > 0$$

The case where  $k > x_q$  is treated similarly.

**(b)** Let  $n$  be **even** and  $q = \frac{n}{2}$ . Then, the median is  $\frac{x_q + x_{q+1}}{2}$ .



We proceed as follows:

(i) for any  $k < x_q$  we prove that

$$\sum_{i=1}^n |x_i - x_q| < \sum_{i=1}^n |x_i - k|$$

(ii) for any  $k > x_{q+1}$  we prove that

$$\sum_{i=1}^n |x_i - x_{q+1}| < \sum_{i=1}^n |x_i - k|$$

(iii) for any  $k, \mu \in [x_q, x_{q+1}]$  we prove that

$$\sum_{i=1}^n |x_i - k| = \sum_{i=1}^n |x_i - \mu|$$

The (i) and (ii) can be proved using the rationale adopted in the case where  $n$  is odd.

For the (iii) case, assuming that  $k < \mu$ , we have

$$B_1 = \sum_{i=1}^n |x_i - k| = \sum_{i=1}^q |x_i - k| + \sum_{i=q+1}^n |x_i - k| = \sum_{i=1}^q (k - x_i) + \sum_{i=q+1}^n (x_i - k)$$

and

$$B_2 = \sum_{i=1}^n |x_i - \mu| = \sum_{i=1}^q |x_i - \mu| + \sum_{i=q+1}^n |x_i - \mu| = \sum_{i=1}^q (\mu - x_i) + \sum_{i=q+1}^n (x_i - \mu)$$

Taking the difference between  $B_1$  and  $B_2$ , we have

$$\begin{aligned} B_1 - B_2 &= \sum_{i=1}^q [(k - x_i) - (\mu - x_i)] + \sum_{i=q+1}^n [(x_i - k) - (x_i - \mu)] = \sum_{i=1}^q (k - \mu) + \sum_{i=q+1}^n (\mu - k) \\ &= q(k - \mu) + (n - q)(\mu - k) = \frac{n}{2}(k - \mu) + \left(n - \frac{n}{2}\right)(\mu - k) = 0 \end{aligned}$$

Therefore, all  $k \in [x_q, x_{q+1}]$  minimize the quantity

$$A = \sum_{i=1}^n |x_i - k|.$$