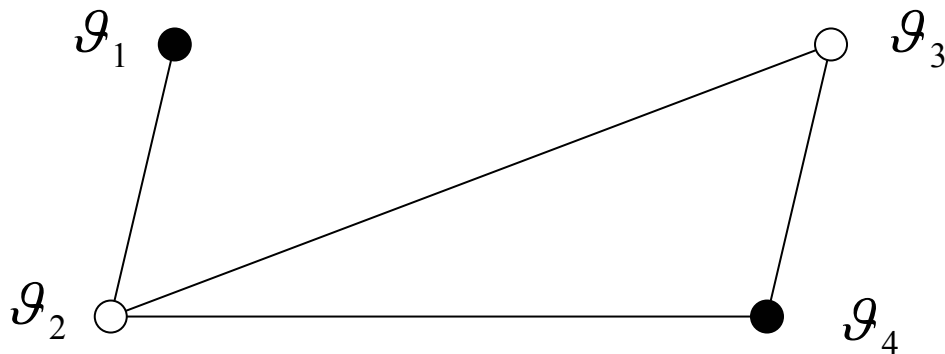


The Local Search: The MIS

example : The maximum Independent Set problem in a graph $G(V,E)$

"Find $V' \subseteq V$ s.t. $\forall u, \mathcal{G} \in V' \Rightarrow [u, \mathcal{G}] \notin E$ "



$\{\mathcal{G}_1, \mathcal{G}_3\}$ a solution



The MIS by the Local Search

Solution coding : $x_1, x_2, \dots, x_i, \dots, x_n$

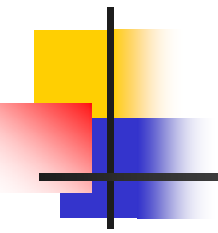
$$x_i = 1 \Rightarrow \mathcal{G}_i \in V'$$

$$x_i = 0 \Rightarrow \mathcal{G}_i \notin V'$$


Function :

$$\text{Max} \sum_{i=1}^n x_i - \lambda |E'|, \lambda \in \mathfrak{R}$$

$$E' = \{[u, \mathcal{G}] \in E \text{ with } u, \mathcal{G} \in V'\}$$



Neighborhood : FLIP

 $x_1, x_2, x_3, \dots, x_i, \dots, x_n$ current sol.
 $x_1, x_2, x_3, \dots, \overline{x_i}, \dots, x_n$ neighbor sol.

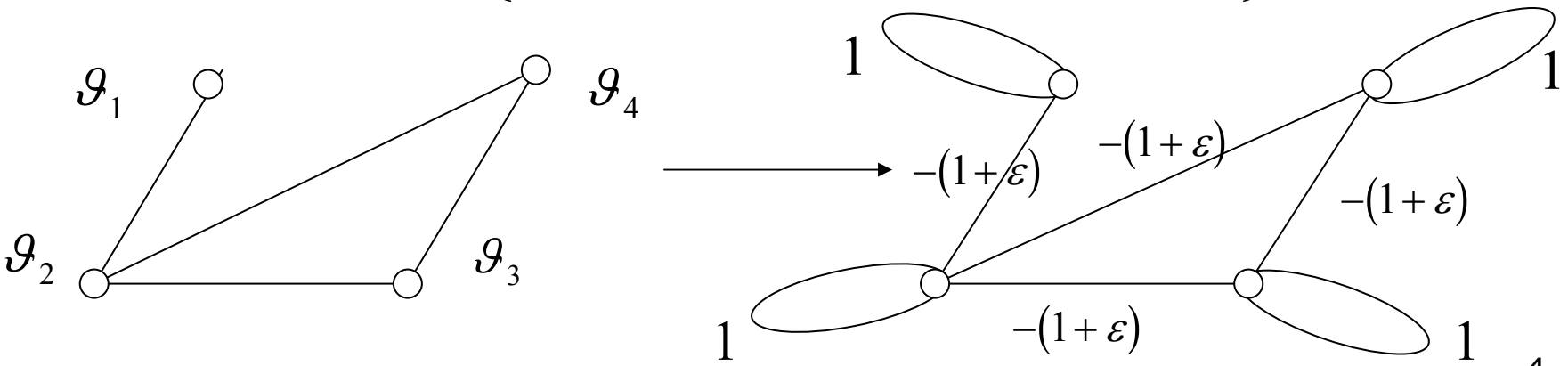
The maximum Independent Set problem

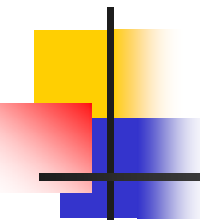
Another solution :

$$G(V, E) \xrightarrow{\text{tranf.}} G'(V, E', W)$$

$$E' = E \cup \{(\mathcal{G}, \mathcal{G}) \mid \mathcal{G} \in V\}$$

$$W = (W_{ij}) = \left\{ \begin{array}{ll} 1 & \text{si } i = j \\ -(1 + \varepsilon) & \text{si } i \neq j, \varepsilon \in \mathfrak{R} \end{array} \right\}$$





Solution coding : n-binary vector of 0,1

$$\mathcal{G} \rightarrow \begin{array}{cccc} \square & \square & \dots & \square \\ \mathcal{G}_1 & \mathcal{G}_2 & & \mathcal{G}_n \end{array}$$

initial configuration : randomly (not feasible)

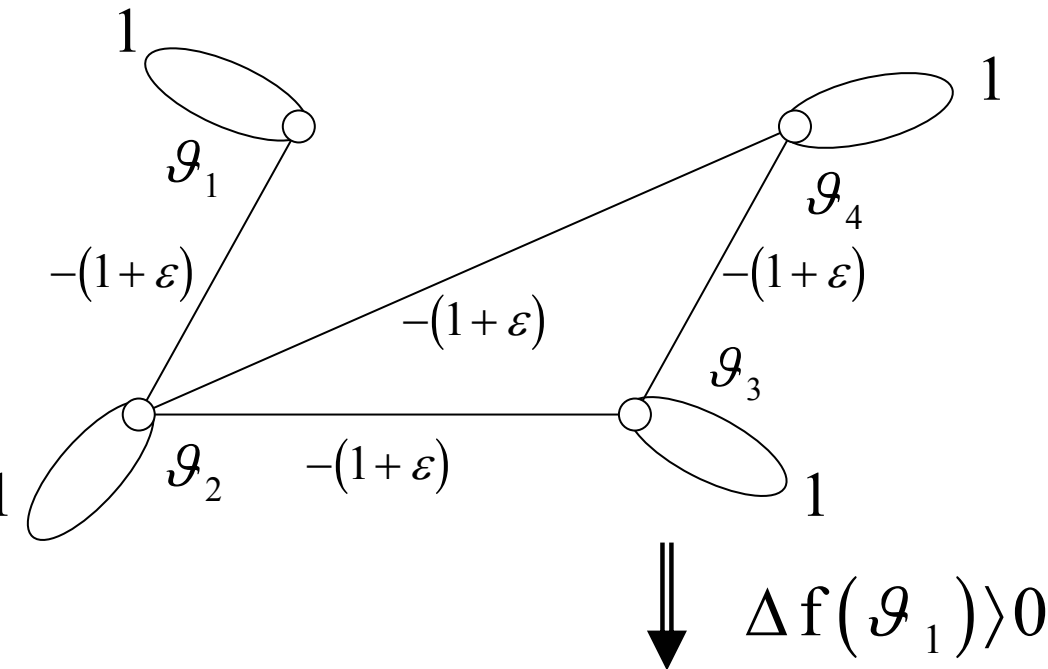
Rule : change bit $\mathcal{G}_i \left\{ \begin{array}{l} \text{yes} \quad \text{if } \Delta f(\mathcal{G}_i) > 0 \\ \text{non} \quad \text{if } \Delta f(\mathcal{G}_i) \leq 0 \end{array} \right\}$

$$\Delta f(\mathcal{G}_i) = (1 - 2\mathcal{G}_i) \left(\sum_{j \in N(\mathcal{G})} w_{ij} \mathcal{G}_j + w_{ii} \right)$$

Example (LS) with the neighborhood FLIP

● \leftrightarrow "1"

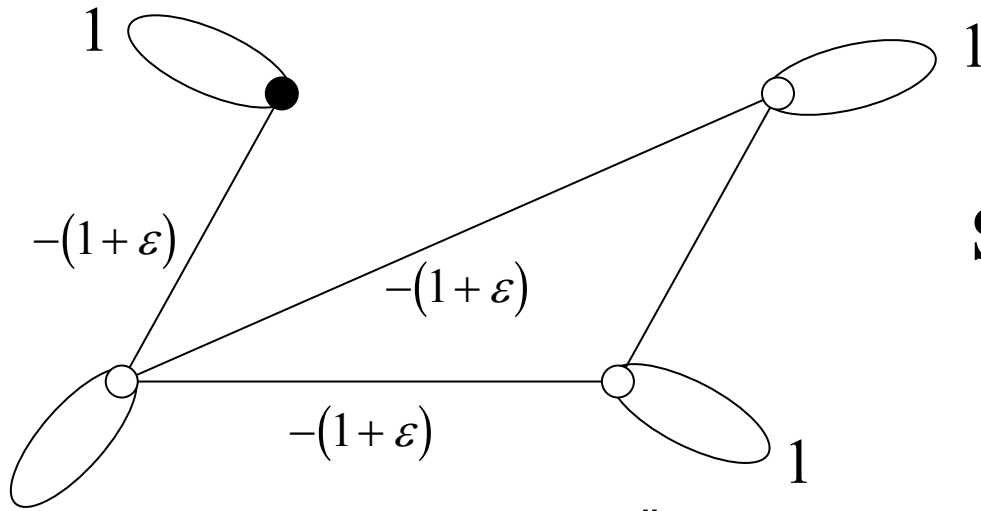
○ \leftrightarrow " \emptyset "



sol. :

\mathcal{G}_1	\mathcal{G}_2	\mathcal{G}_3	\mathcal{G}_4
○	○	○	○

MIS: Another Solution (the second vertex is happy)

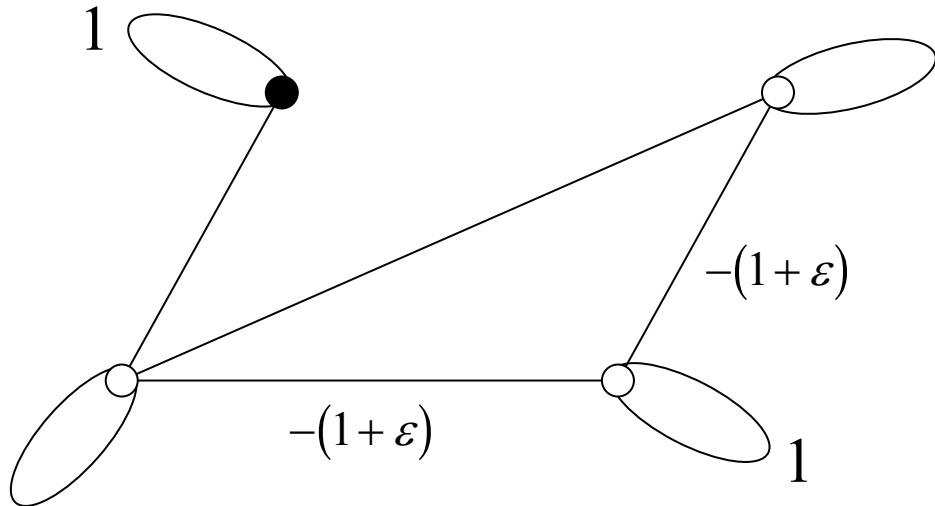


sol. : 1 0 0 0



$$\Delta f(\mathcal{G}_2) < 0$$

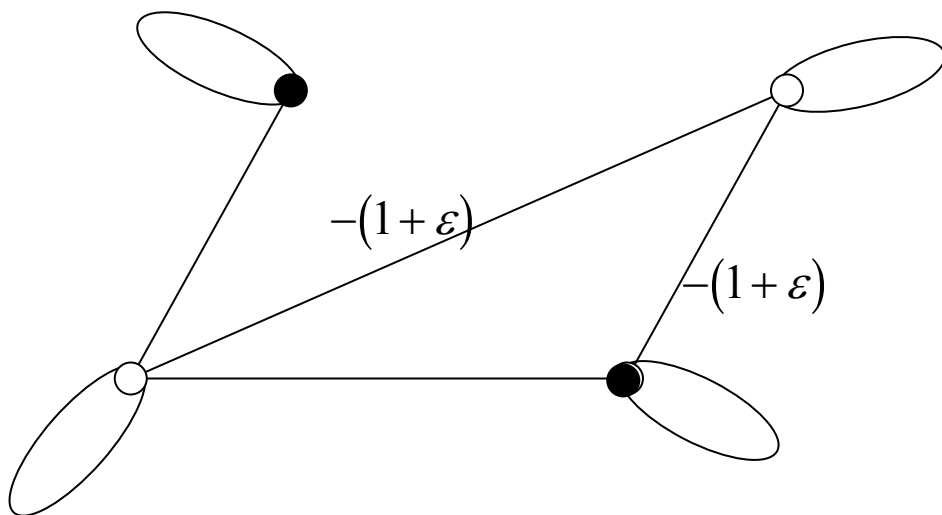
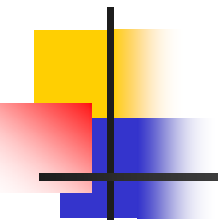
The third vertex becomes Happy



sol. : 1 0 0 0



$$\Delta f(\mathcal{G}_3) > 0$$



sol. : 1 0 1 0
system
" HAPPY "



$$\Delta f(\mathcal{G}_4) < 0$$

$$\Delta f(\mathcal{G}_1) < 0$$

•

•



Rermarks

- each vertex \longrightarrow "happy" locally
- vertices \longrightarrow "happy" asynchronously
- each vertex \longrightarrow simple operations (addition , comp.)

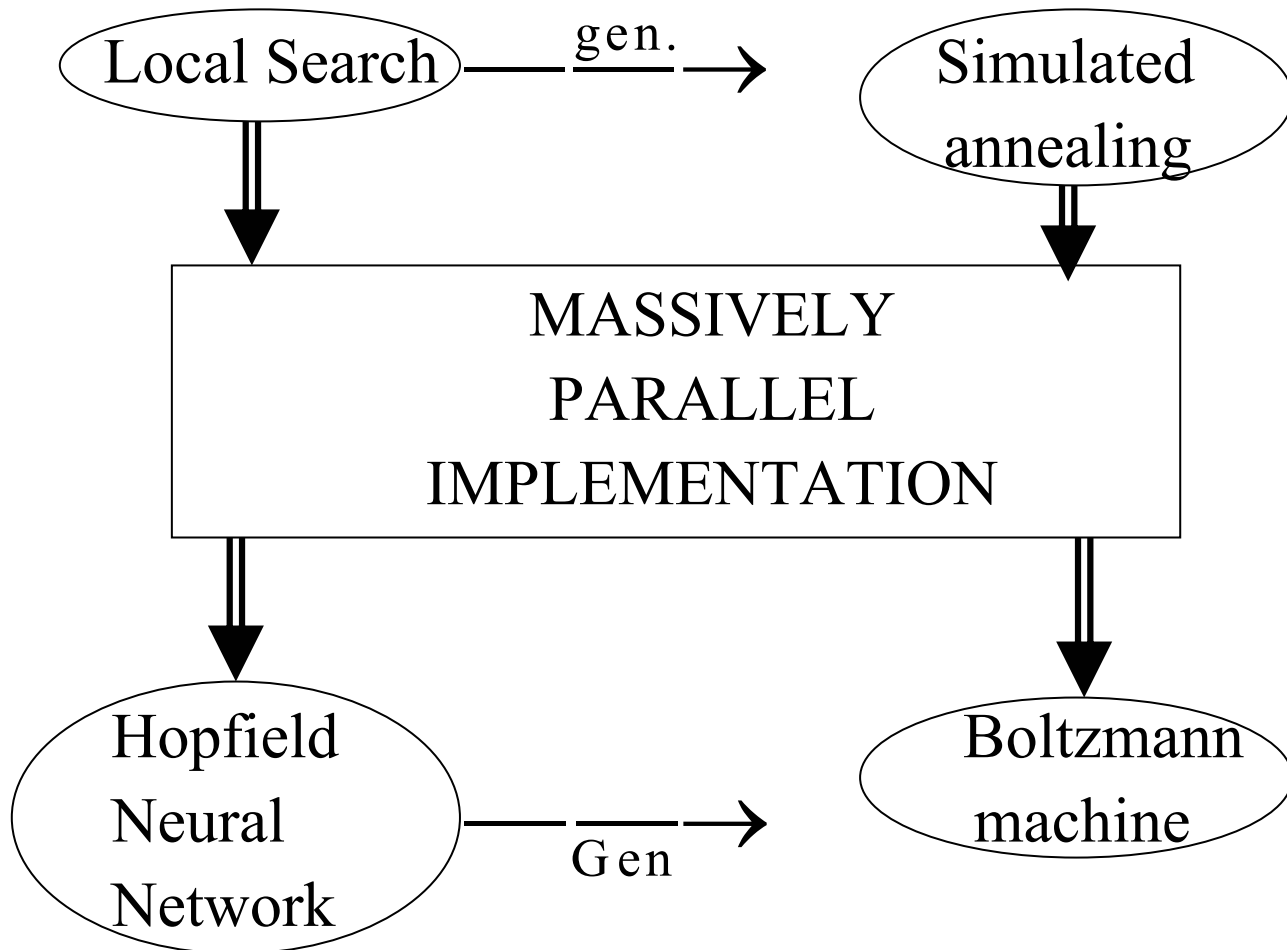
"Massively Parallel algorithm" (Hopfield model)

- + Simulated annealing



Asymptotically optimal algorithm
(Boltzmann machine)

CONCLUSION





Future Work

- A problems classification relatively to HARDNESS with meta-heuristics
 - Neighborhoods for using in practice !
 - Landscapes RUGGEDNESS
 - ... A THEORY ... → ?